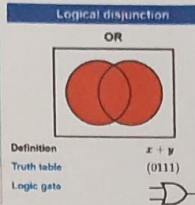


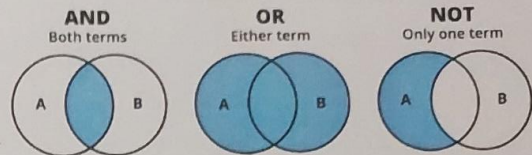
Logical addition
(disjunction)

A	B	F=A ∨ B
0	0	0
0	1	1
1	0	1
1	1	1

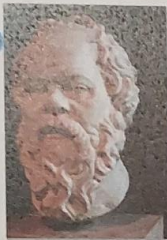
A	B	A ∨ B
True	True	True
True	False	True
False	True	True
False	False	False



BOOLEAN LOGIC

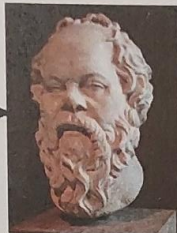
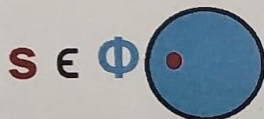


Good logic



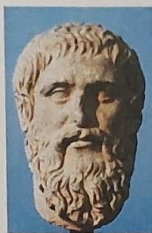
Socrates

Socrates was
a philosopher

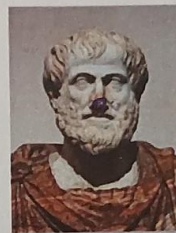


Socrates

philosophers are men



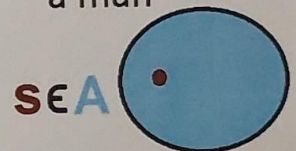
Plato



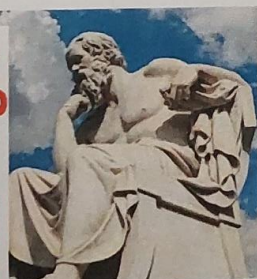
Aristotle



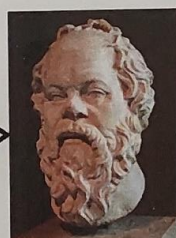
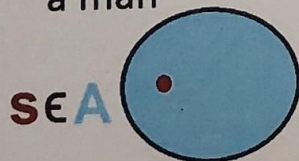
Socrates was
a man



Bad logic

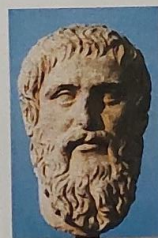
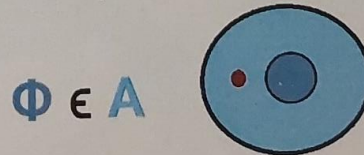


Socrates was
a man

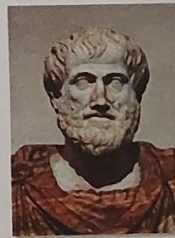


Socrates

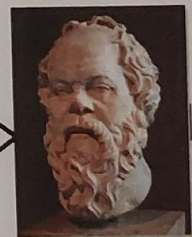
philosophers are men



Plato

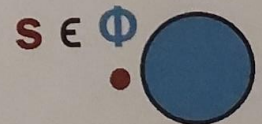


Aristotle



Socrates

Socrates was
a philosopher



Resume of Lecture by Pr. Bob Gallagher from MIT

MIT Massachusetts Institute of Technology (MIT)

George Boole (1815-1864) developed Boolean logic

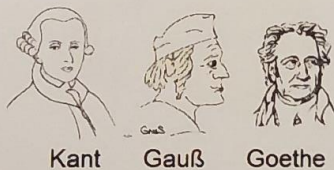
The principles of logical thinking have been understood (and occasionally used) since the Hellenic era.

Boole's contribution was to show how to systemize these principles and express them in equations (called Boolean logic or Boolean algebra).

Claude Shannon (1916-2001) showed how to use Boolean algebra as the basis for switching technology. This contribution systemized logical thinking for computer and communication systems, both for the design and programming of the systems and their applications.

Logic continues to be abused in politics, religion, and most non-scientific areas.

Logic continues to be abused in politics, religion and most non-scientific areas



Kant

Gauß

Goethe

A little nationalistic, but this is an sample of right logic

Kant, Gauss, Goethe are great

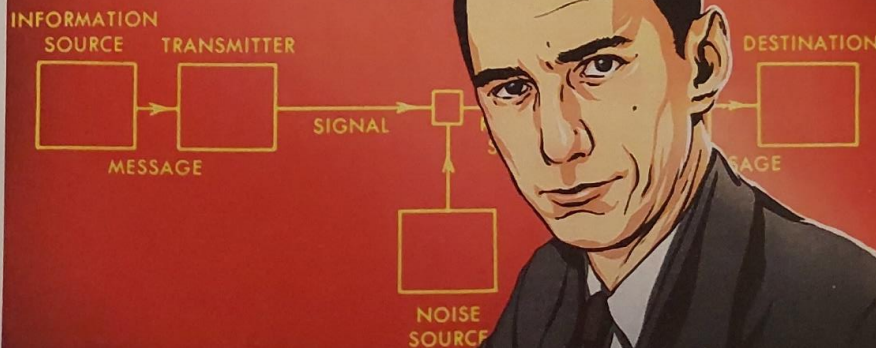
Kant, Gauss, Goethe - Germans

Germany Great



Bad logic (abuse of logic)

The Mathematical Theory of Communication



Creating a reliable connection over an unreliable (noisy) channel that's what IT is about

and that's what Shannon did

Fri 16th 11° 7°
 Walking in Oxford on a cold and rainy day
 With prof. Matthias Winkel
 DEPARTMENT OF STATISTICS OXFORD



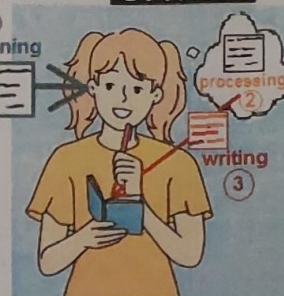
80% chance of rain

says the Met Office in its weather forecast for Oxford.

CHALK+TALK

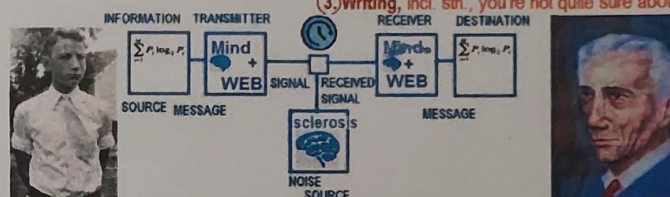


ink + think



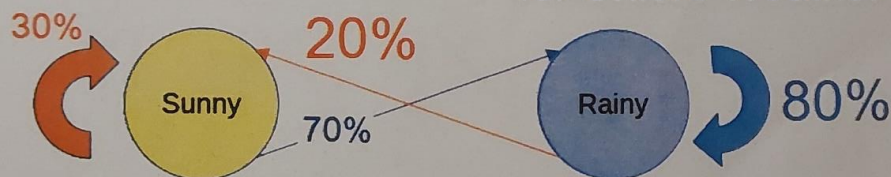
Good way ① listening ② first way of processing

③ Writing, incl. sth. you're not quite sure about

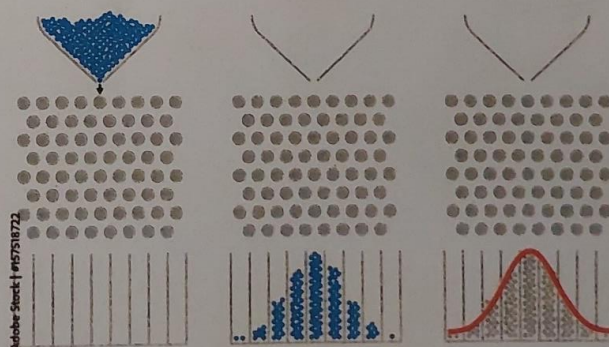
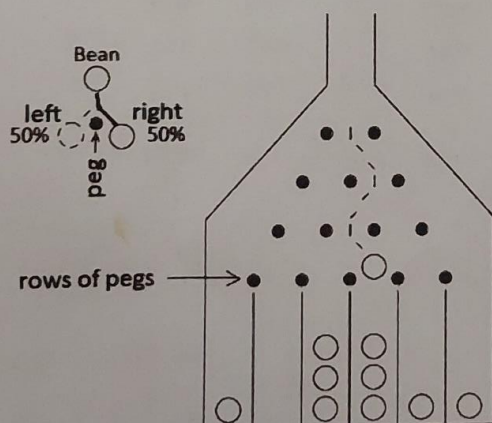


Markoff Chain Probability Model

for Oxford Weather



If it is Rainy today => there is an 80% chance that it will be rainy tomorrow.



School $\downarrow$ gravity $\downarrow$ MOTION == formalism ==> University $E = MC^2$ $\#515.6$ $W = 2uf$ $\| \vec{J} \vec{a} \vec{s}$

Motivation: 80% chance of rain

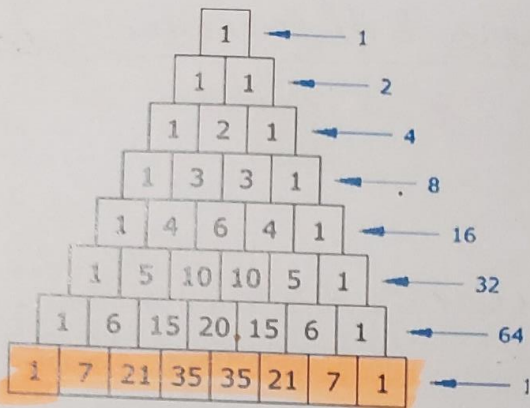
Let A_j be the event of rain at 9am on day j of this term, $1 \leq j \leq n$

Suppose the events A_i each have probability p , independently

Oxford			
Tue 13th	Wed 14th	Thu 15th	Fri 16th
10° 9°	13° 10°	13° 8°	11° 7°
70%	70%	70%	80%

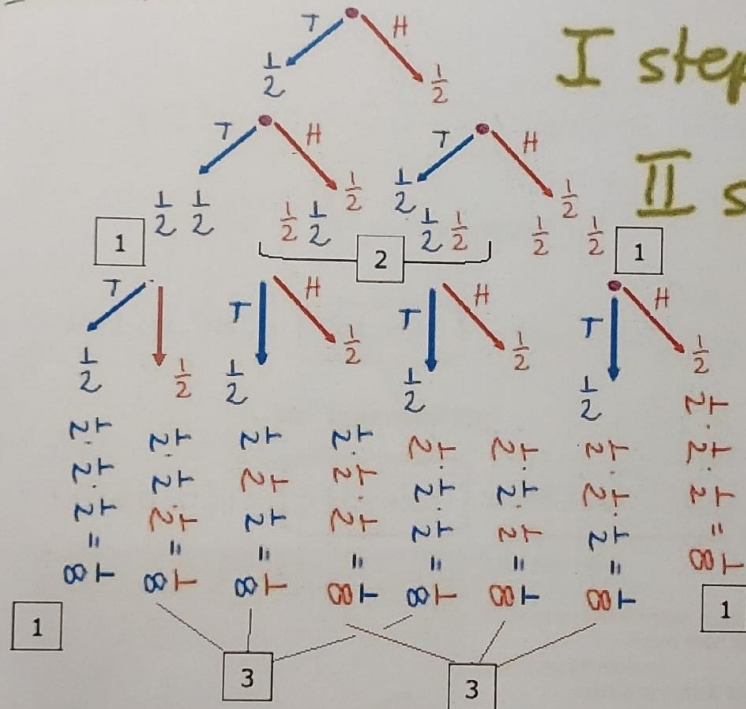


then take notes on the lecture yourself



Pascal's Triangle (Binomial coefficients)

8 row 1 8 28 56 70 56 28 8 1



Shannon Hartley Theorem

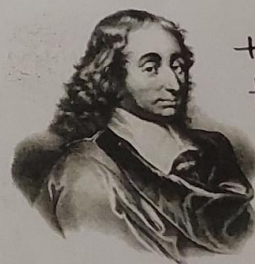
This is a measure of the *capacity* on a channel; it is impossible to transmit information at a faster rate without error.

$$C = B \log_2(1 + S/N)$$

- C = capacity (in bit/s)
- B = bandwidth of channel
- S = signal power (in W)
- N = noise power (in W)

It is more usual to use SNR (in dB) instead of power ratio (as with terrestrial and commercial communications systems).
S/N >> 1, then rewriting in terms of log10.

$$C = B \frac{\log_{10}(S/N)}{\log_{10} 2} = B \frac{10 \log_{10}(S/N)}{10 \cdot \log_{10} 2} = B \frac{SNR}{3.01}$$



+ 0.2
to Exam
from B.
Pascal

$$\begin{aligned} (a+b)^0 &= 1 \\ (a+b)^1 &= a+b \\ (a+b)^2 &= a^2 + 2ab + b^2 \\ (a+b)^3 &= a^3 + 3a^2b + 3ab^2 + b^3 \\ (a+b)^4 &= a^4 + 4a^3b + 6a^2b^2 + 4ab^3 + b^4 \\ (a+b)^5 &= a^5 + 5a^4b + 10a^3b^2 + 10a^2b^3 + 5ab^4 + b^5 \end{aligned}$$

$$(1+x)^8 = 1^8 + 8x + 28x^2 + 56x^3 + 70x^4 + 56x^5 + 28x^6 + 8x^7 + x^8$$

Generalization of
The number of arrangements

$$\underbrace{x_1 \dots x_1}_{m_1} \underbrace{x_2 \dots x_2}_{m_2} \dots \underbrace{x_k \dots x_k}_{m_k}$$

$$m_1 + m_2 + \dots + m_k = n$$

$$\frac{n!}{m_1! m_2! \dots m_k!}$$

$$C_m^n = \frac{n!}{m!(n-m)!} = \binom{n}{m}$$

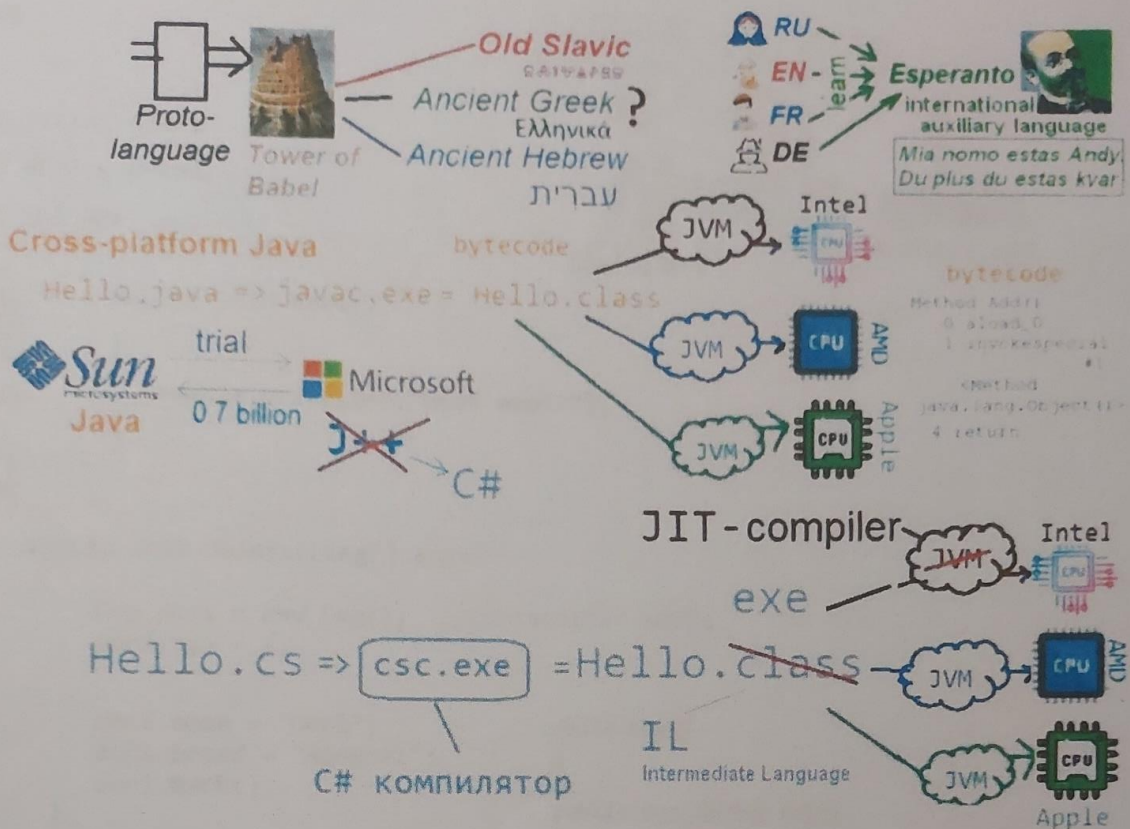
$$(1+x)^5 = 1^5 + 5x + 10x^2 + 10x^3 + 5x^4 + x^5$$

$$(a+b)^7 = a^7 + 7a^6b + 21a^5b^2 + 35a^4b^3 + 35a^3b^4 + 21a^2b^5 + 7ab^6 + b^7$$

symmetrical

Binomial coefficients: $x^0 + 4x^1 + 6x^2 + 4x^3 + x^4$

$$\binom{4}{0} \binom{4}{1} \binom{4}{2} \binom{4}{3} \binom{4}{4}$$



Hello.vb ⇒ **vbc.exe** = ~~Hello.class~~ .exe
 VB компилятор

Hello.pl ⇒ **plc.exe** = ~~Hello.class~~ .exe

CIL - Common Intermediate Language

```

void .ctor() {
    .maxstack 1
    IL_0000: ldarg.0
    IL_0001: call instance
    void
    System.Object::.ctor()
    IL_0006: ret
} // end of method Add::ctor
  
```



After him I love
 More than I love these eyes,
 more than my life,
 More, by all mores,
 than e'er I shall love wife



Windows Kernel

JMP -CorEXEMain
CIL (named .EXE)

```
class Dog
```

```
{
```

```
    public string name;
```

```
    public string breed;
```

```
    public int age;
```

```
    public void Bark()
```

```
    {
```

```
        Console.WriteLine("Woof woof!");
```

```
    }
```

```
static void Main(string[] args)
```

```
{
```

```
    Dog dori = new Dog(); //Constructor works
    dori.age = 3;
```

```
    dori.name = "Dori";
```

```
    dori.breed = "Mongrel";
```

```
    dori.Bark();
```

```
}
```

1 - **constructor** no returns value

2. The name of the **constructor** is the same name is the class.

3. more than one **constructor**

```
public Dog()
```

```
{
```

```
}
```

```
public Dog(string name)
```

```
{
```

```
    this.name = name;
```

```
}
```

```
public Dog(string name, string breed)
```

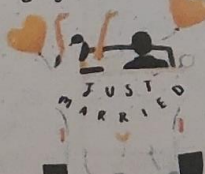
```
{
```

```
    this.name = name;
```

```
    this.breed = breed;
```

```
}
```

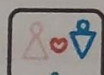
Congratulations



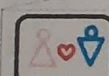
private protected internal



private



protected



internal



public

```
class AFather
```

```
{
```

```
    protected string name;
```

```
    int age;
```

```
}
```

```
class ASon:AFather
```

```
{
```

```
    public ASon(string name)
```

```
    {
```

```
        base.age = 33;
```

```
        base.name = name;
```

```
    }
```

```
}
```

```
class Program
```

```
{
```

```
    static void Main(string[] args)
```

```
    {
```

```
        AFather af = new AFather();
```

```
        ASon andy = new ASon("Olaf");
```

```
    }
```

```
}
```

unavailable available

Program.exe

```
class AFather
```

```
{
```

```
    protected string name;
```

```
    internal int age;
```

```
}
```

```
class Program
```

```
{
```

```
    static void Main(string[] args)
```

```
    {
```

```
        AFather af = new AFather();
```

```
        af.age = 33;
```

```
        ASon andy = new ASon("Olaf");
```

```
    }
```

```
}
```

ASon.dll

```
class ASon:AFather
```

```
{
```

```
    public ASon(string name)
```

```
    {
```

```
        base.name = name;
```

```
        base.age = 33;
```

```
    }
```

```
}
```



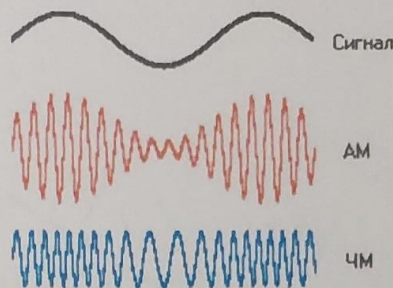
internal



protected



Reginald A. Fessenden
(October 6, 1866 – July 22, 1932)

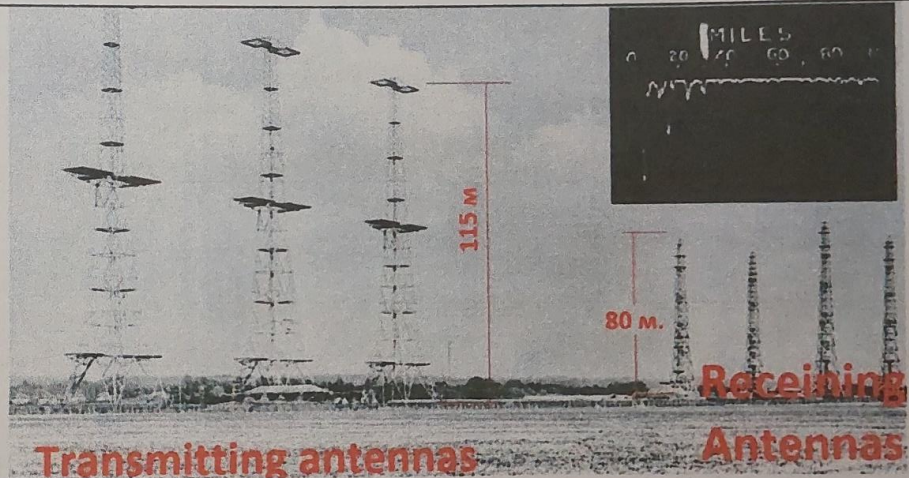


(October 6, 1866 – July 22, 1932)

first transmission of speech by radio (1900), and the first two-way radiotelegraphic communication across the Atlantic Ocean (1906)

"Ни одна организация, занимающаяся какой-либо конкретной областью деятельности, никогда не изобретает какие-либо важные разработки в этой области или не внедряет какие-либо важные разработки в этой области до тех пор, пока она не будет вынуждена сделать это из-за внешней конкуренции.." Oxford University Press. The Quarterly Journal of Economics , Feb., 1926, p. 262.

Battle of Britain
(3 month 3 weeks)
10.07-31.10.1940



Radar played a major role in the Battle of England

H. Nyquist



$$W = K \log m$$

Where W is the speed of transmission of intelligence,
 m is the number of current values,
and, K is a constant.

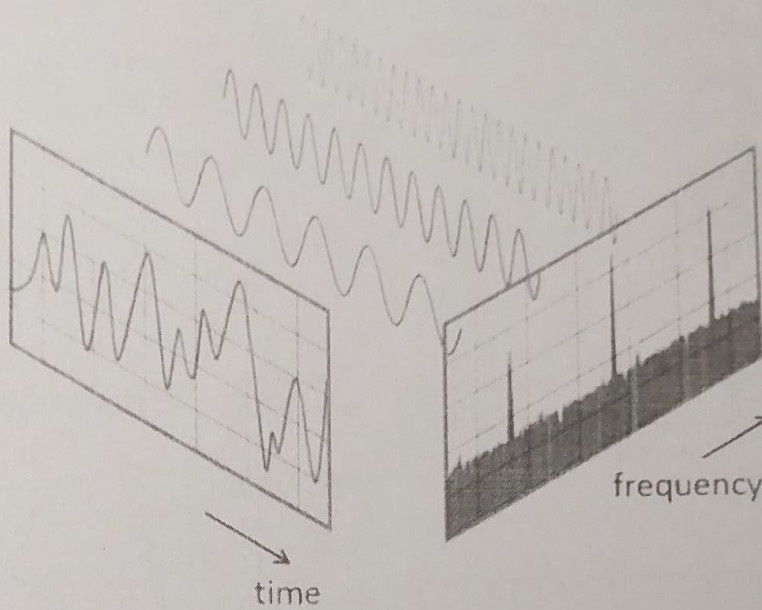
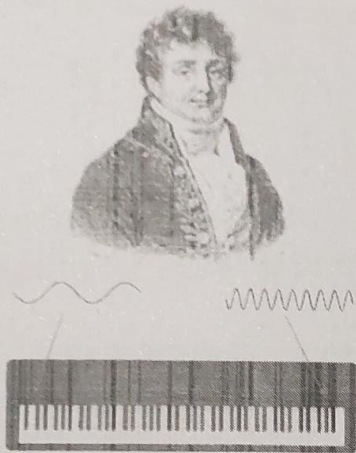


Ralph **Hartley**
(81:1888-1970)

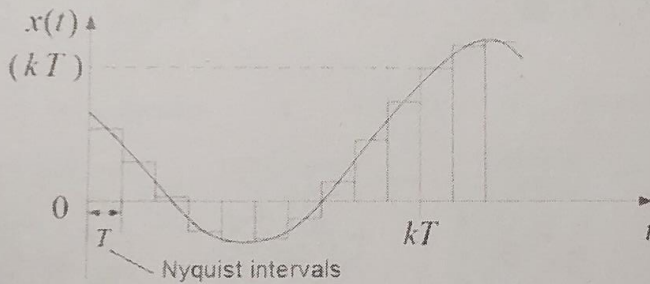
$$H = n \log s$$

$$= \log s^n.$$

Fourier transform



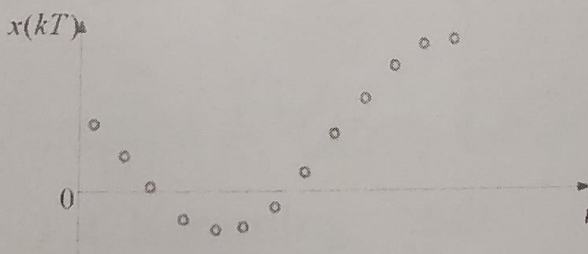
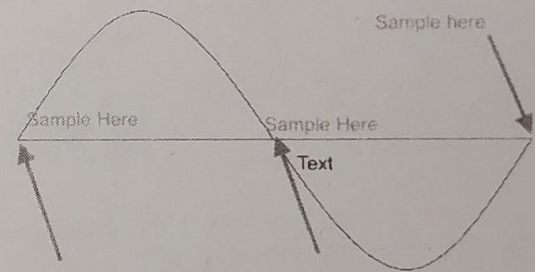
Sampling. Kotelnikov-Nyquist Theorem



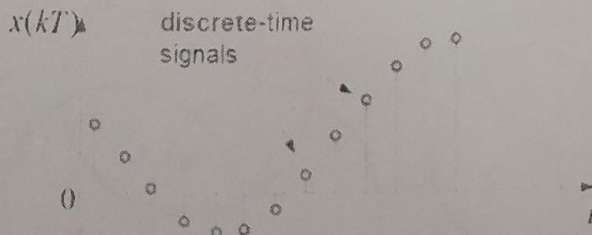
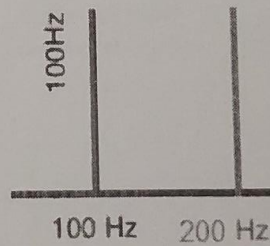
Time intervals T , through which readings $s(kT)$ are taken, are called Nyquist intervals.

Sine with period T

Sampling at $T/2$



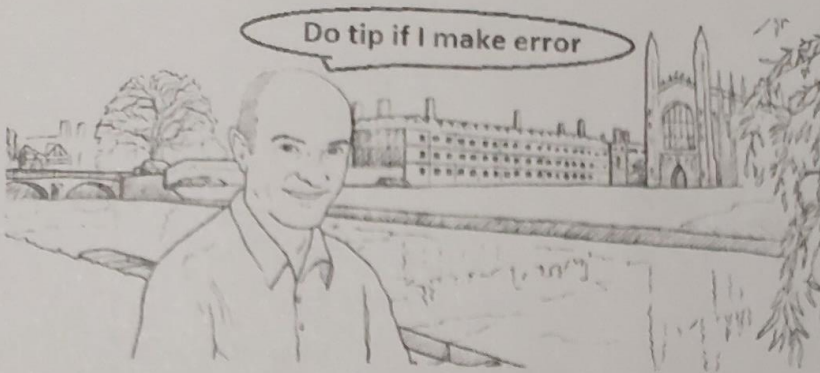
frequency Sample



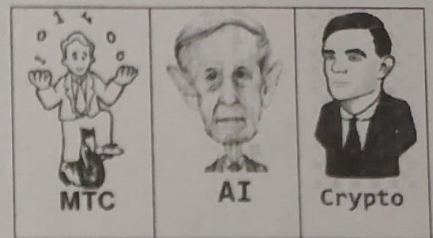
$$F_{\text{sample}} \geq 2 * F_{\text{max}}$$

$$(T_{\text{sample}} \leq T_{\text{min}} / 2)$$

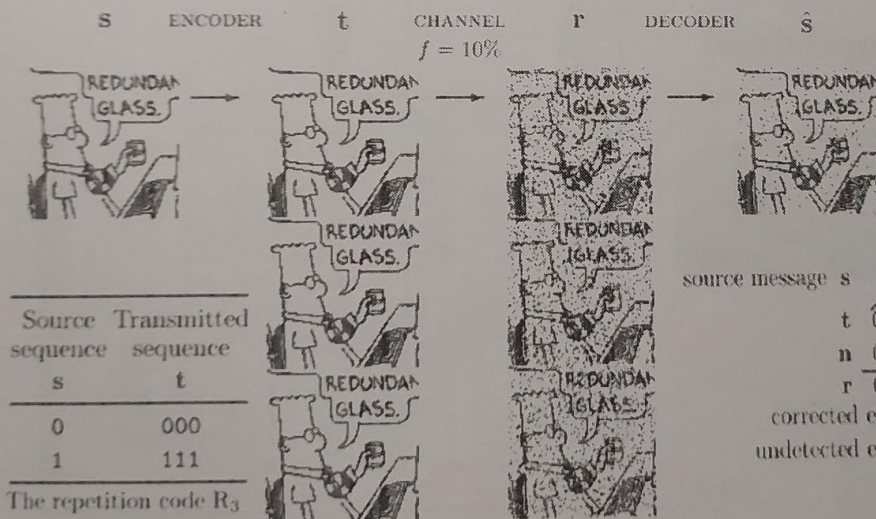
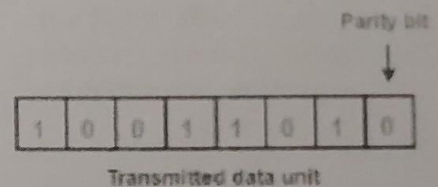
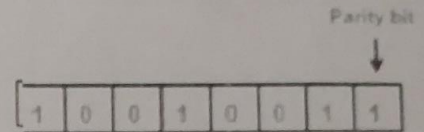
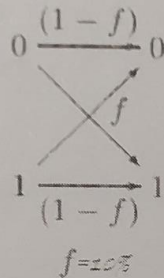
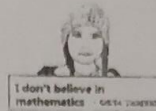
Котельников



Sir Dr. D. MacKay,
University of Cambridge
(22 April 1967 – 14 April 2016)

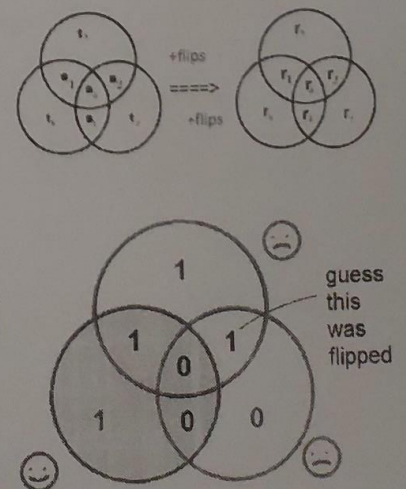
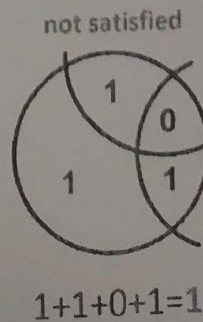
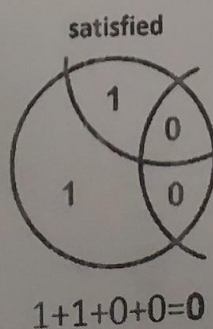
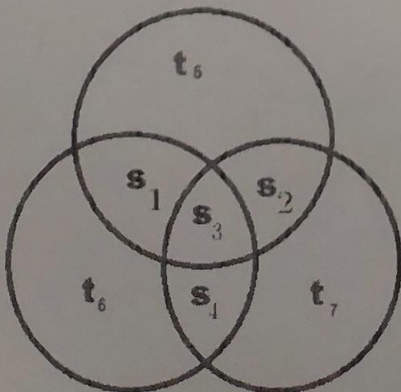


"I believe in clean energy,
but I also believe in mathematics"



source message s	0	0	1	0	1	1	0
t	000	000	111	000	111	111	000
n	000	001	000	000	101	000	000
r	000	001	111	000	010	111	000
corrected errors		*					
undetected errors						*	

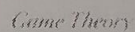
7.4. Hamming code. $\frac{4}{\Sigma} \rightarrow \frac{7}{t}$



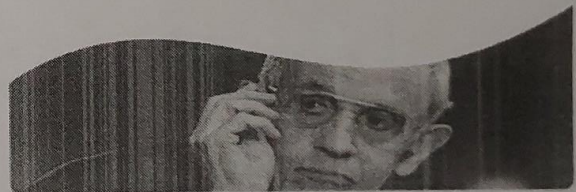


by Vilfredo Pareto

The orange sector E-D-E is the most Pareto efficient - since an increase in one indicator leads to a decrease in another.



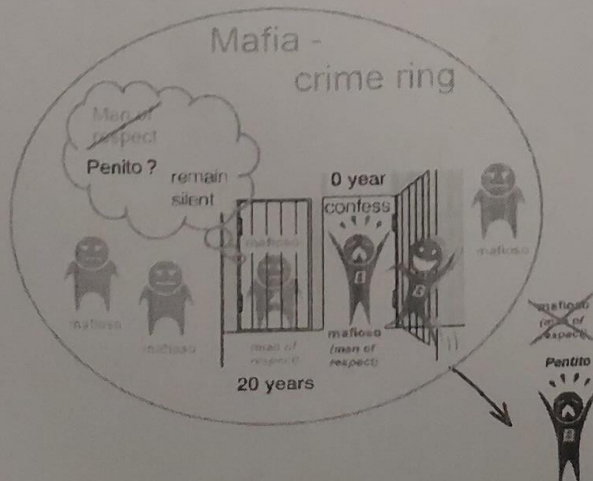
Nash Equilibrium



**** => Nash equilibrium**

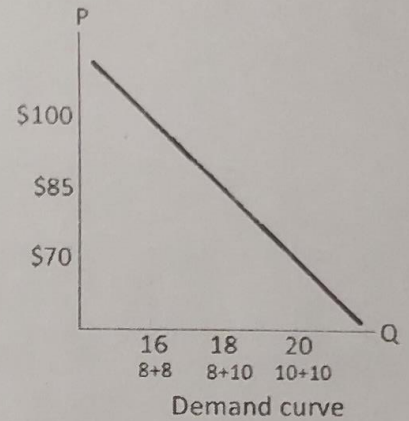
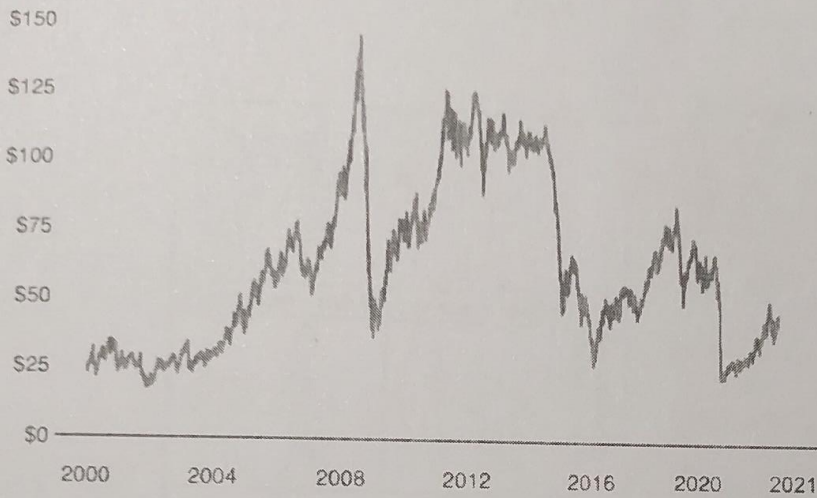
	$H_2(x)$	Recognition;	Non-recognition;
$H_1(x)$	<i>Player 2</i>		
<i>Player 1</i>		1	2
Recognition;	1	-5 * -5 → 0	-20
Non-recognition;	2	-20 0 * ↓ -1	-1

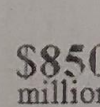
Pareto Optimality



Oil price hits 18-year low

Brent crude, US dollars per barrel



Barrel		1.		2.	
		$8 \cdot 10^6$  day		$10 \cdot 10^6$  day	
1.	$8 \cdot 10^6$ 	 \$800 millions per day  \$100  \$800 millions per day	 \$850 millions per day  \$85  \$680		
	$10 \cdot 10^6$ 	 \$680 millions per day  \$85  \$850 millions per day	 \$700 millions per day  \$70  \$700 millions per day		

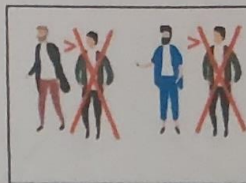




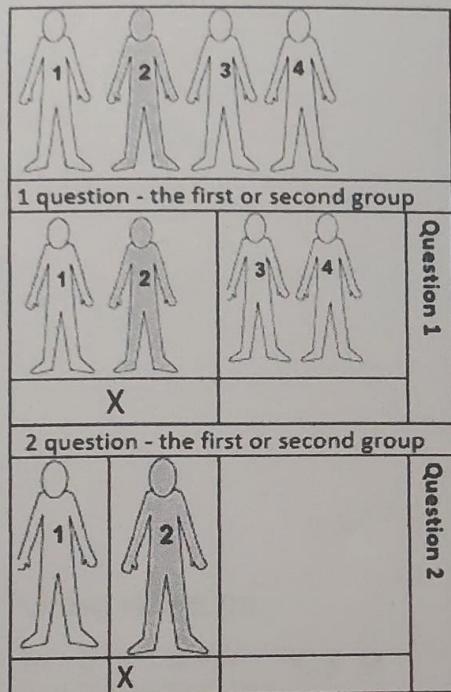
Say **NO** to the first



Say **YES** to the second if it is better than the first



Say **NO** to the third only if it is worse than all the others



Average number of questions =

$$2 \cdot 0.25 + 2 \cdot 0.25 + 2 \cdot 0.25 + 2 \cdot 0.25 = 2$$

Average number of questions =

$1 \cdot 0.5 +$	$2 \cdot 0.25 +$	$3 \cdot 0.125 +$	$3 \cdot 0.125$
			

Question 1. Is this Zuckerberg?	 50%	$1 \cdot 0.5$
Question 2. Is this Sergey Brin?	 25%	$2 \cdot 0.25$
Question 3. Is this Stefan from BMW?	 12,5%	$3 \cdot 0.125$
So Prince Saud	 12,5%	$3 \cdot 0.125$
Average number of questions = 1,75		

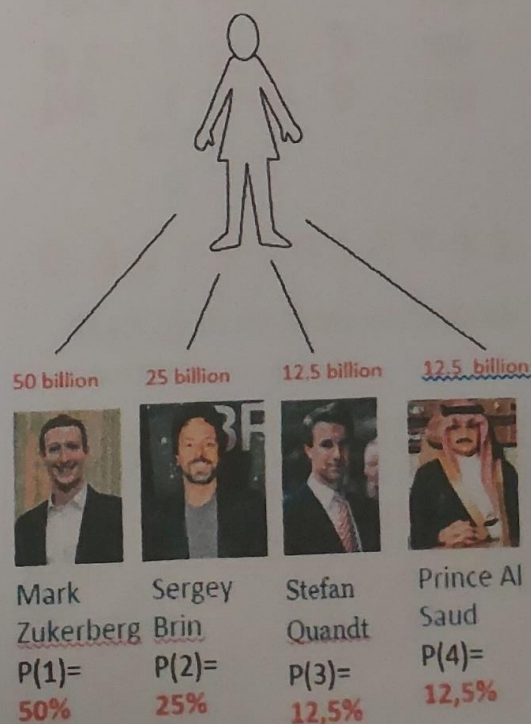
Quantifying information

$$S(x) = \sum_{i=1}^n p(i) \log_2 \frac{1}{p(i)}$$

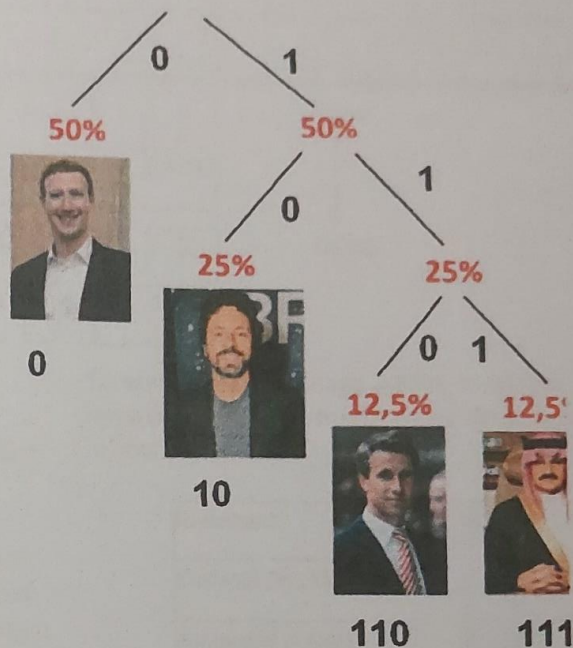
Quantifying information

$$I(x_i) = \log_2 \left(\frac{1}{p_i} \right)$$

number of bits required to encode choice



Huffman code



First-order approximation
(symbols independent but with frequencies of Belarusian txt).

Мама мыла ра	
М - 3 — 30%	1-3 М
а - 4 — 40%	4-7 а
ы - 1 — 10%	8 -ы
л - 1 — 10%	9 -л
р - 1 — 10%	10 -р
10	
лла мама р	



Мама мыла ра

Ма - 2 22%	1-2 ма
ам - 2 22%	3-4 ам
мы - 1 11%	5 мы
ыл - 1 11%	6 ыл
ла - 1 11%	7 ла
ар - 1 11%	8 ар
ра - 1 11%	9 ра

9

0. 4 6 7 3 1 9 1 6 7 3 5
 ам ыл ла ам ма ра ма ыл ла ам мы
 мылла рама



Second-order approximation (digram (2-symbols) structure as in Belarusian)



Caesar Cipher

$$C = (p + 3) \bmod 26$$

- We can use the ordinal positions of letters in a cipher to generate this key:
- We can also rotate the starting point. If we add 3 to every number, we might use this key:

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C

A	B	C	D	...	Y	Z
1	2	3	4	...	25	26

A	B	C	D	...	Y	Z
4	5	6	7	...	2	3

ABBA

DEED

Vigenere Cipher

An improvement we can make to the Caesar cipher is to increase the number of keys.



While the Caesar cipher uses a single key, the **Vigenere** cipher uses multiple keys by selecting a keyword.

In the Vigenere cipher, for each new letter of message, it is enciphered using a different

letter of the keyword.

<https://www.youtube.com/watch?v=BgFJD7oCmDE>

To encrypt the message **ABBA** using the keyword **LAW**, we might come up with the following table:

Plaintext	A	B	B	A
Ordinal Position	1	2	2	1
Keyword (LAW)	L	A	W	L
Keyword Ordinal Position	12	1	23	12
Sum	13	3	25	13
Ciphertext	M	C	Y	M

Frequency Analysis <https://www.youtube.com/watch?v=sMOZf4GN3oc> Khan

- Another issue with Caesar ciphers is that an adversary may be able to crack the code without a pin.
- For example, if we see a single letter word in the message, we might be able to guess that the character or number represents I or A. From there, we might be able to discover some patterns in the message.
- A pattern may be how frequently letters appear in the English language.

A	B	C	D	E	F	G	H	I	J	K	L	M
8.1%	1.5%	2.8%	4.3%	12.7%	2.2%	2.0%	6.1%	7.0%	0.2%	0.8%	4.0%	2.4%
N	O	P	Q	R	S	T	U	V	W	X	Y	Z
6.7%	7.5%	1.9%	0.1%	6.0%	6.3%	9.1%	2.8%	1.0%	2.4%	0.2%	2.0%	0.1%

- Some letters appear very frequently, such as E or T and some letters appear very infrequently, such as J or K. Using these frequencies, we can look at what appears frequently or infrequently in the cipher-text and perhaps find certain patterns.
- While for humans it might be tedious to conduct frequency analysis to decode a message, a computer can do it very quickly.